

## Sommario

- Definizione di ricorsione e strategie *divide et impera*
- Semplici algoritmi ricorsivi
- Merge Sort
- Quicksort
- Esempi più complessi di algoritmi ricorsivi

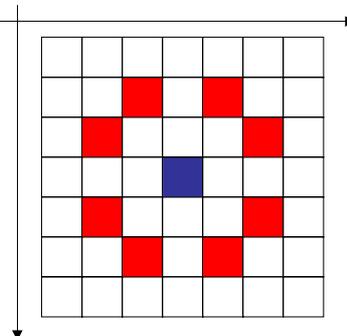
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## Il tour del cavaliere

Si desidera trovare una sequenza di mosse del cavallo tale per cui questo tocchi una ed una sola volta ciascuna casella di una scacchiera  $N \times N$ .

Si ricorda che un cavallo in posizione  $(i,j)$  può muovere in 8 possibili caselle:



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a.a. 2001/2002



## Inizializzazioni

```
#define DIM 6
int  a[8],b[8],scacc[DIM][DIM];
void main(void)
{ int    i,    j,    result;
  a[0]=2; b[0]=1; a[1]=1; b[1]=2;
  a[2]=-1; b[2]=2; a[3]=-2; b[3]=1;
  a[4]=-2; b[4]=-1; a[5]=-1; b[5]=-2;
  a[6]=1; b[6]=-2; a[7]=2; b[7]=-1;

  for( i=0; i<DIM; i++)
    for( j=0; j<DIM; j++)
      scacc[i][j] = 0;
```

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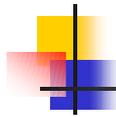


## Programma principale

```
scacc[0][0] = 1;
result = muovi( 2, 0, 0);
if( result == 1)
{ for( i=0; i<DIM; i++)
  { for( j=0; j<DIM; j++)
    printf( "%2d ",scacc[i][j]);
    printf( "\n");
  }
} else
{ printf( "Soluzione non trovata\n");
}
}
```

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a.a. 2001/2002



## Muovi (1)

```
int muovi( int mossa, int posx, int posy)
{  int    i,ret,newposx,newposy;
  if( mossa == (DIM*DIM+1))
    return(1);
  for( i=0; i<8; i++)
  {  newposx = posx + a[i];
    newposy = posy + b[i];
```

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## Muovi (2)

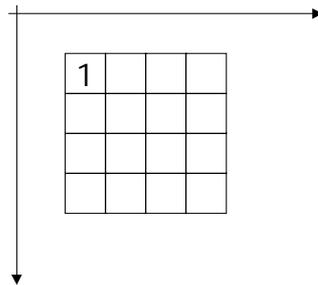
```
  if( (newposx<DIM) && (newposx>=0) &&
      (newposy<DIM) && (newposy>=0))
  {  if( scacc[newposx][newposy] == 0)
    {  scacc[newposx][newposy]=mossa;
      ret=muovi(mossa+1,newposx,newposy);
      if( ret == 0)
        scacc[newposx][newposy]=0;
      else
        return(1);
    }
  }
  return(0);
}
```

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# Analisi

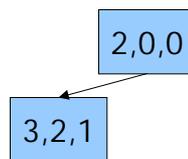
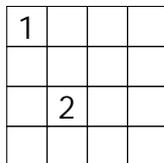
Si supponga che  $N=4$ .



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# Mossa 2

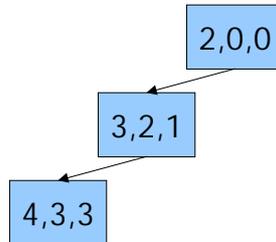


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## Mossa 3

1			
	2		
			3

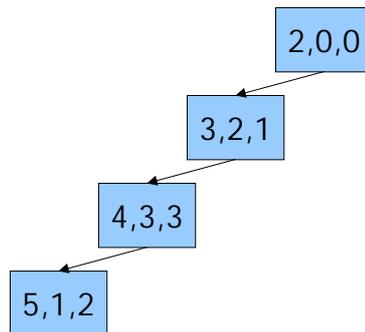


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## Mossa 4

1			
		4	
	2		
			3

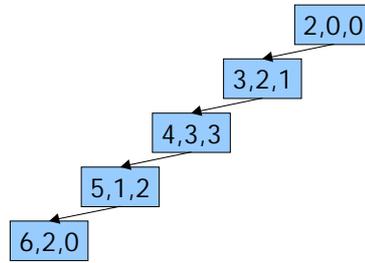


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## Mossa 5

1			
		4	
5	2		
			3

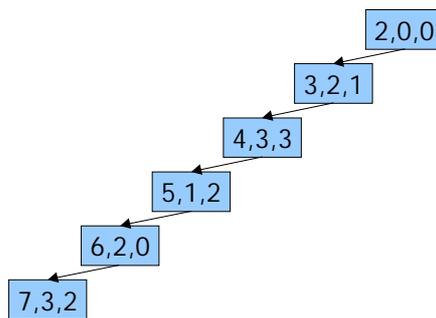


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## Mossa 6

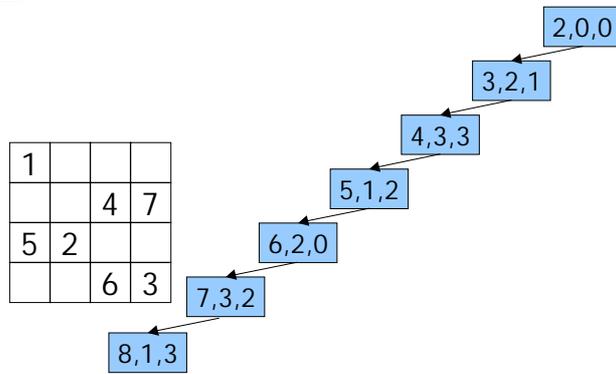
1			
		4	
5	2		
		6	3



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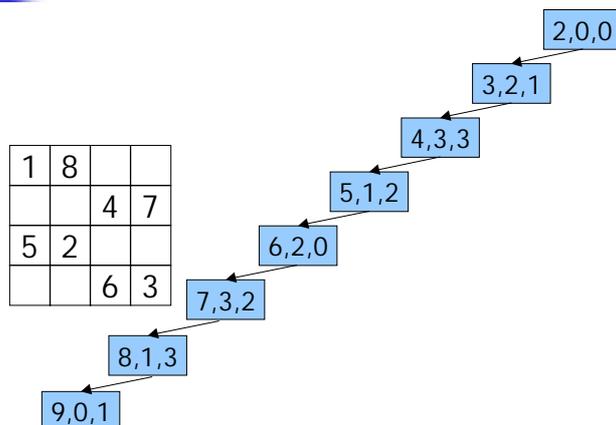
## Mossa 7



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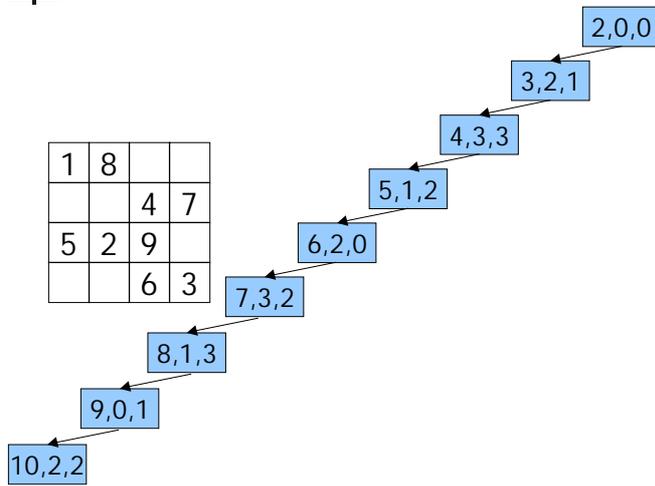
## Mossa 8



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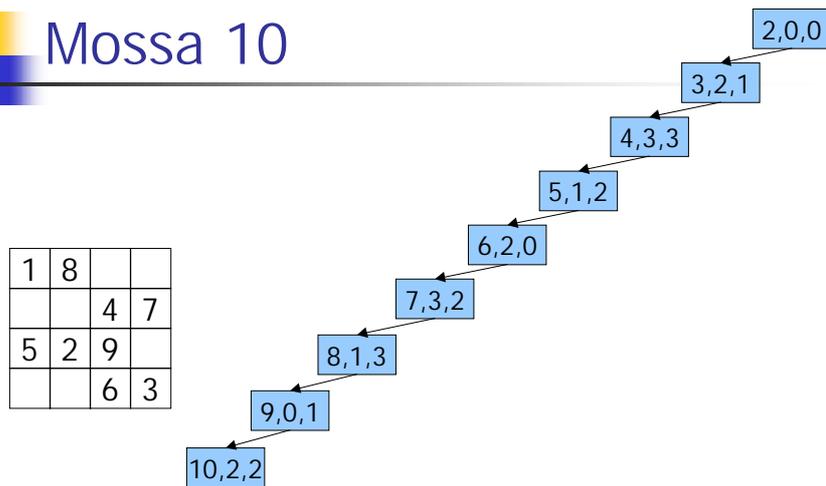
## Mossa 9



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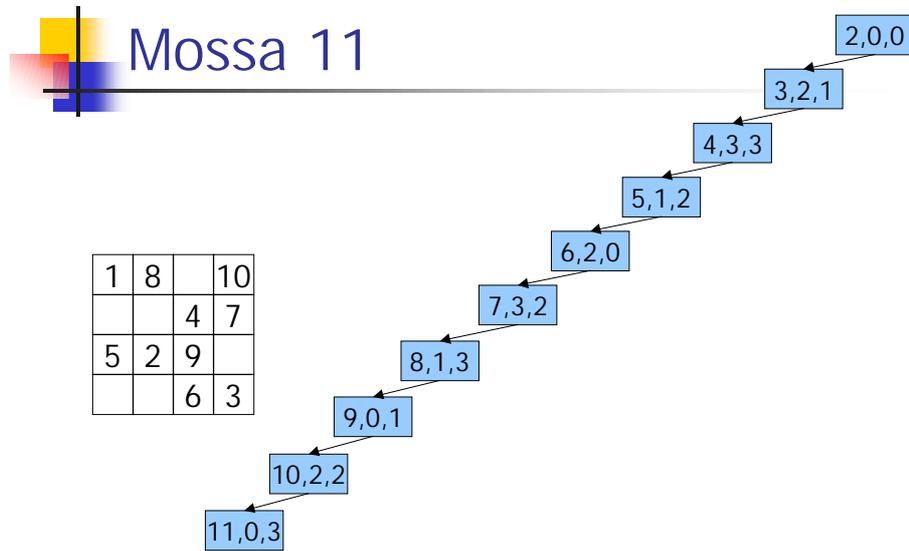
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## Mossa 10



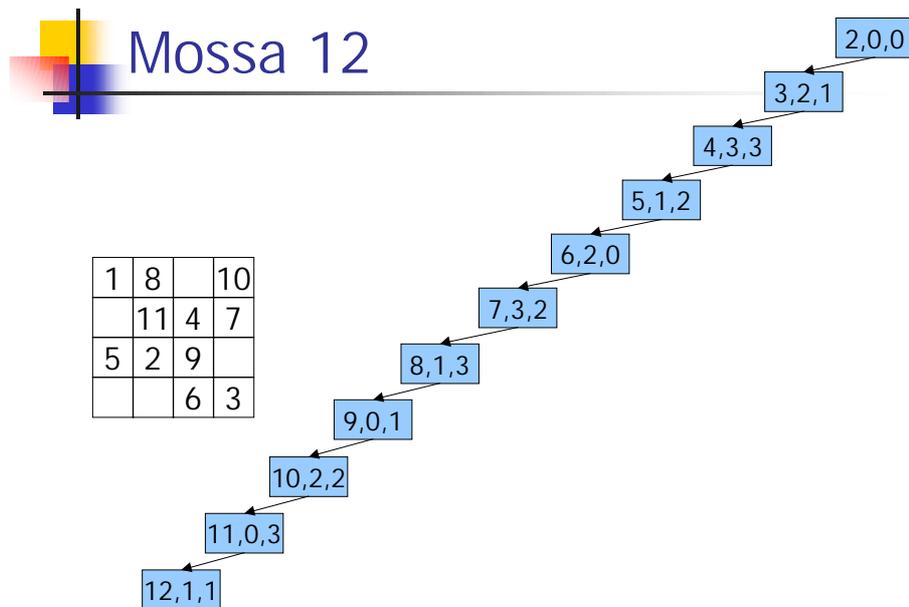
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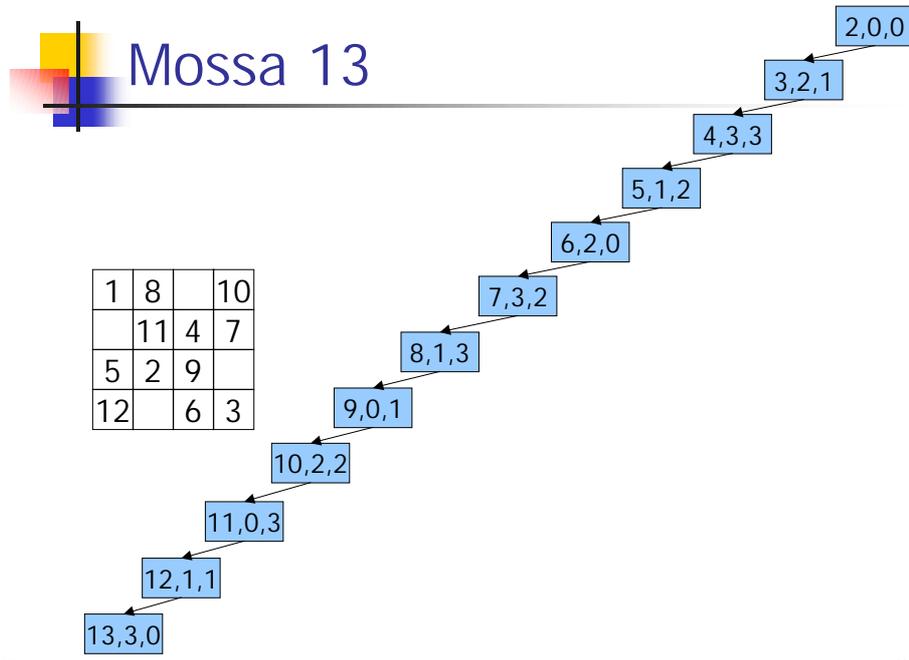
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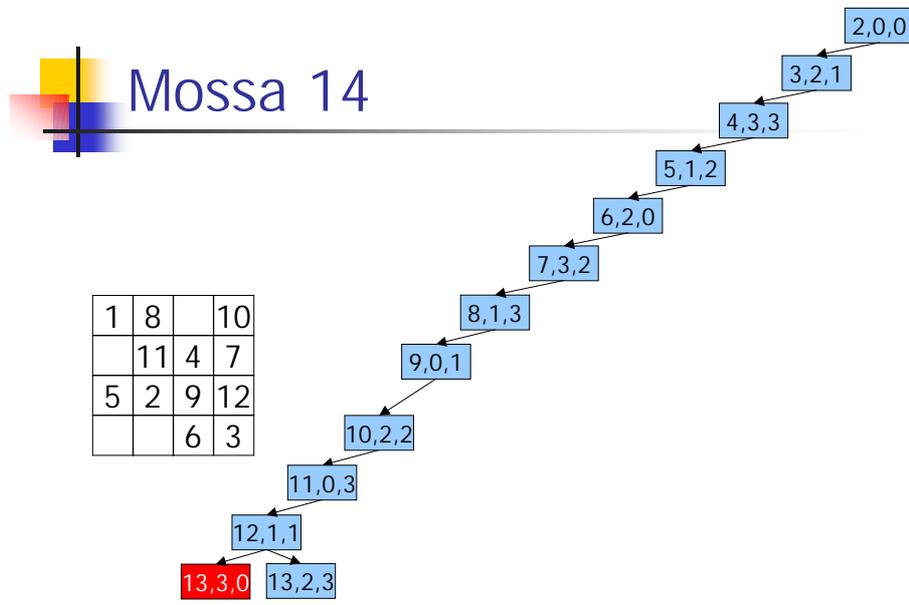
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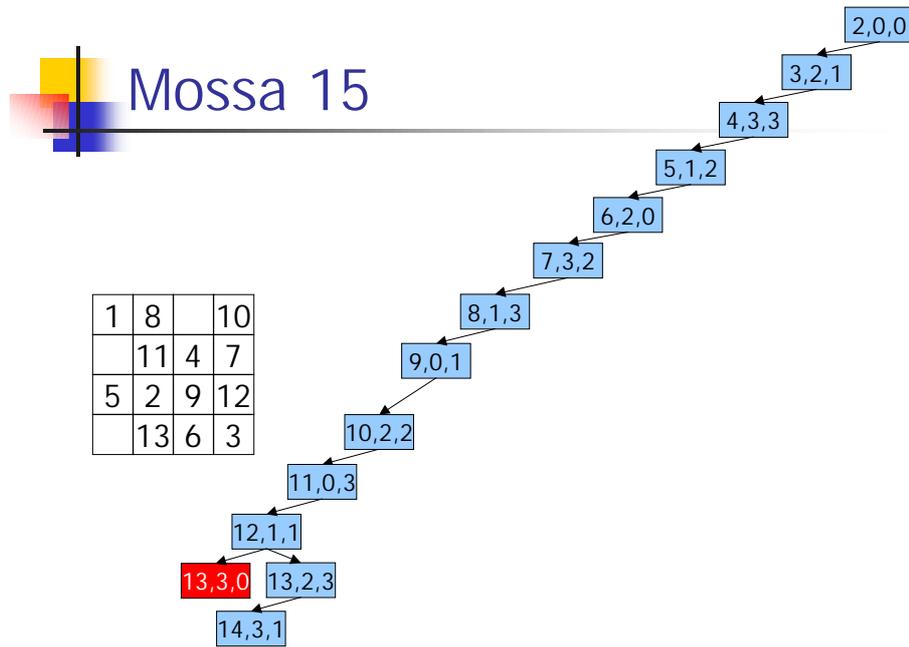
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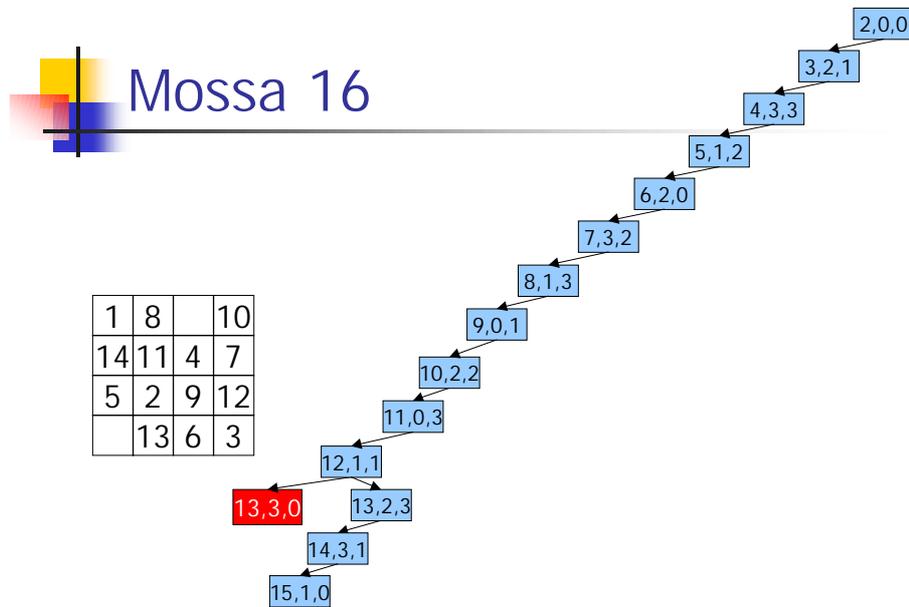
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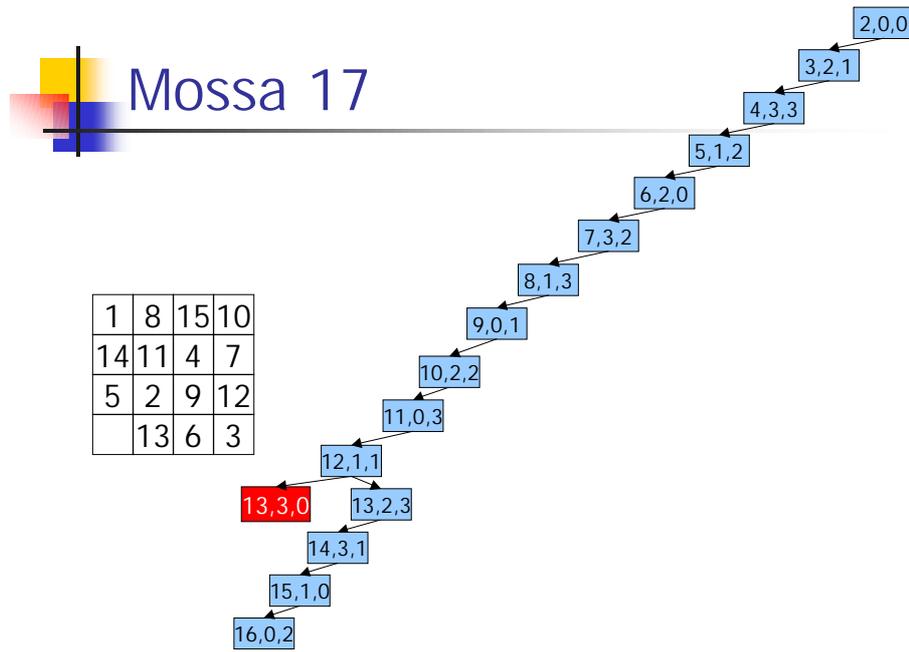
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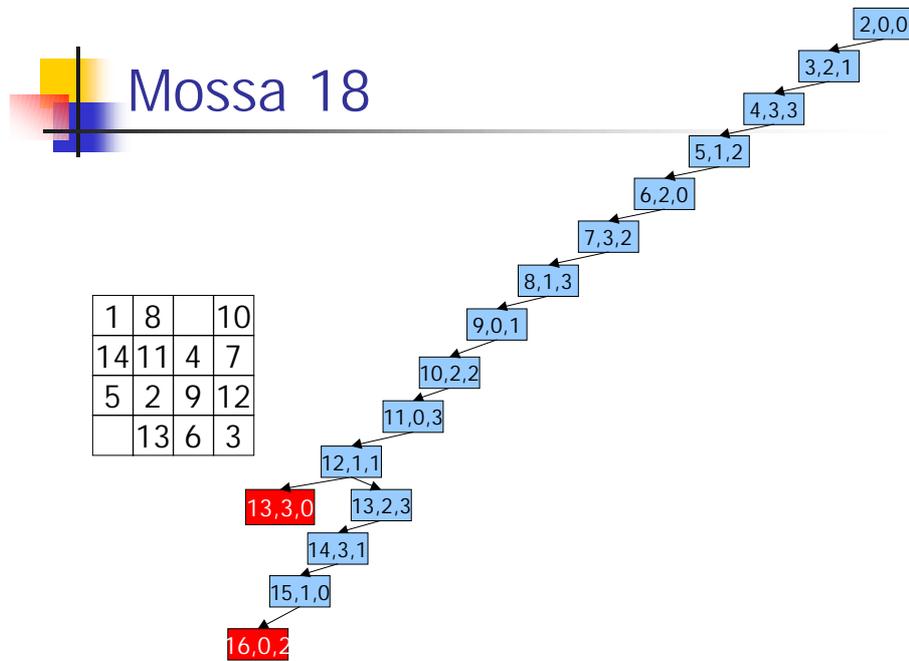
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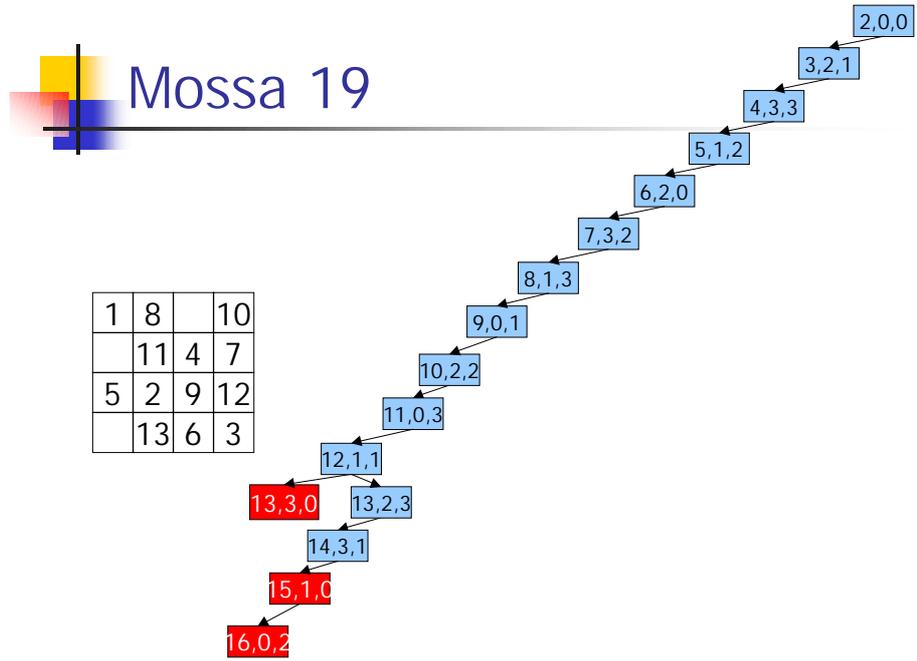
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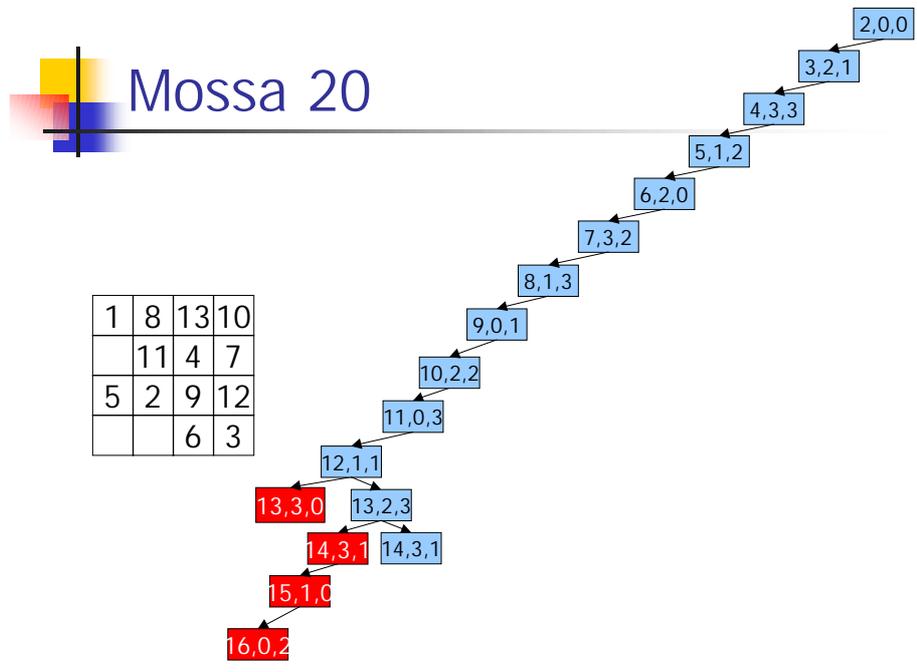
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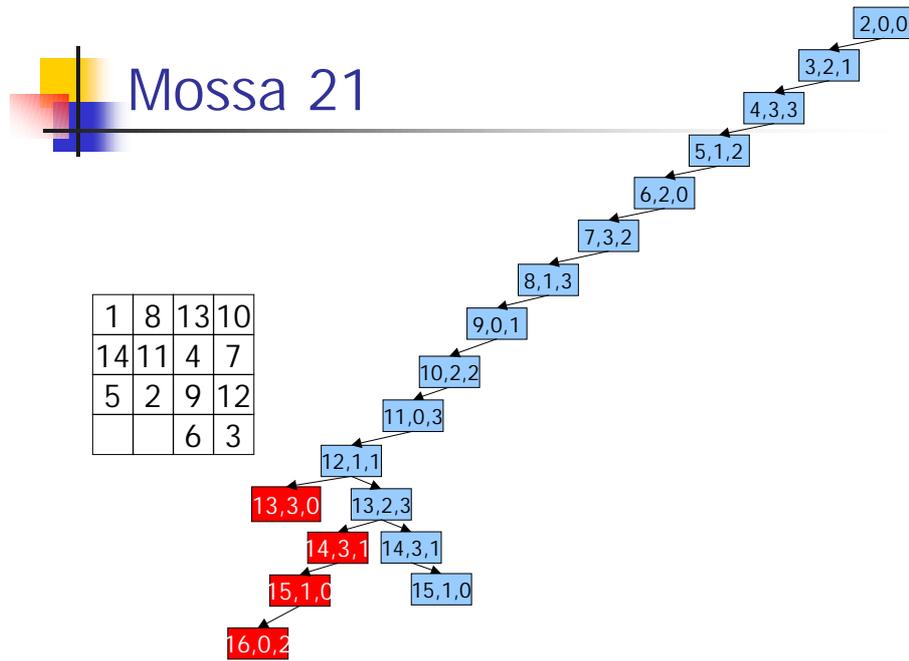
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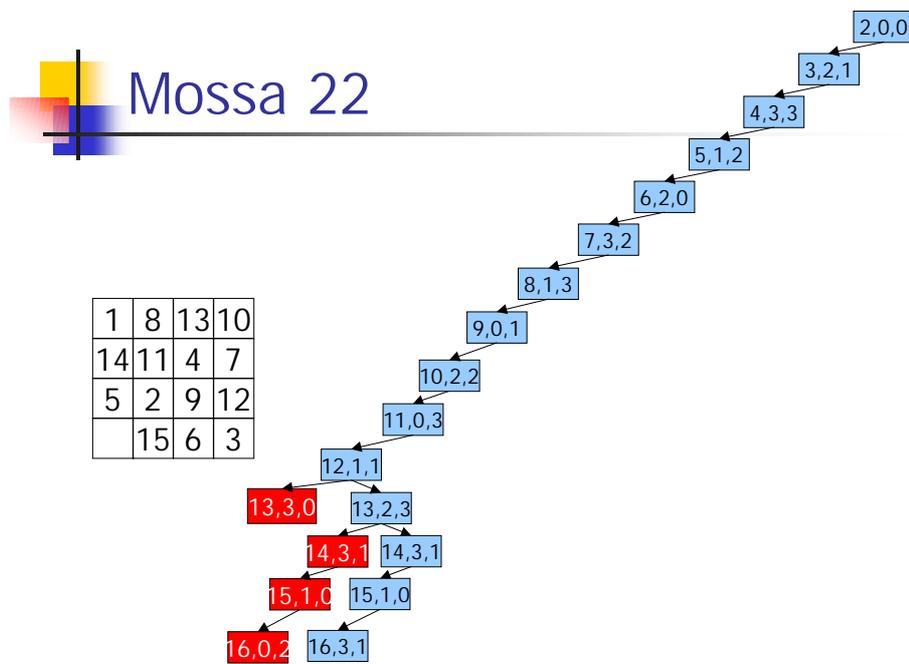
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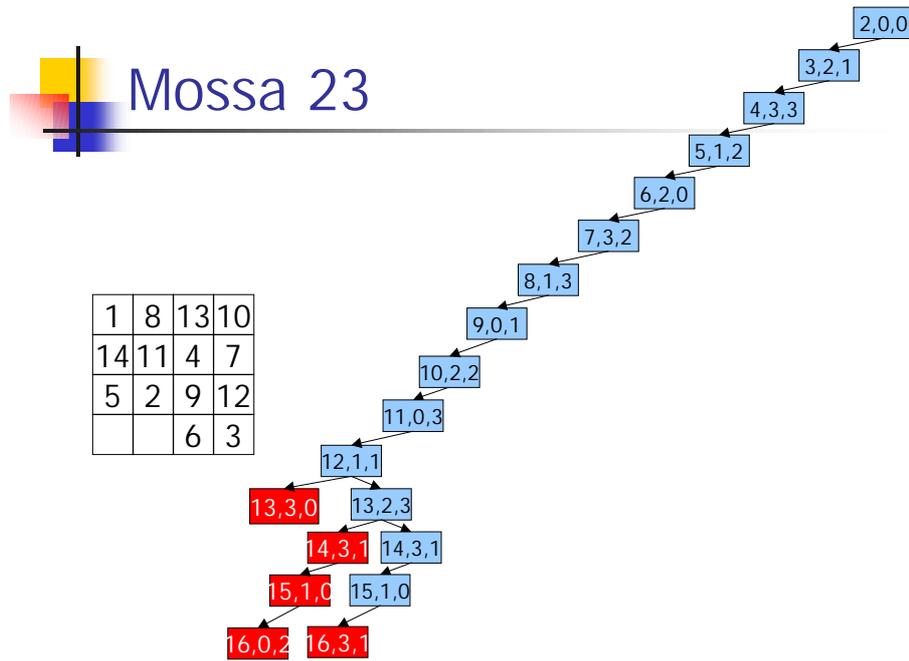
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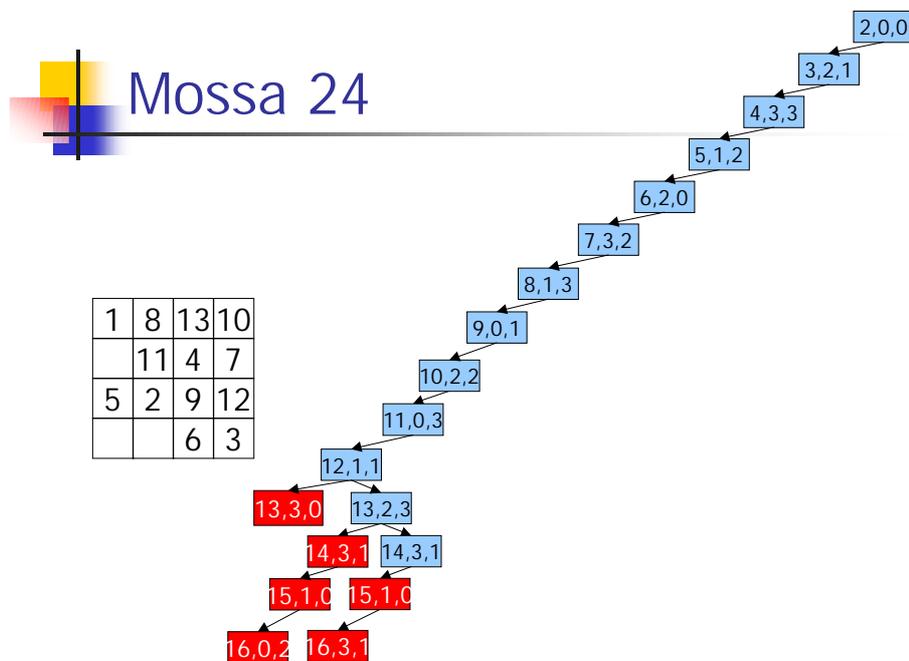
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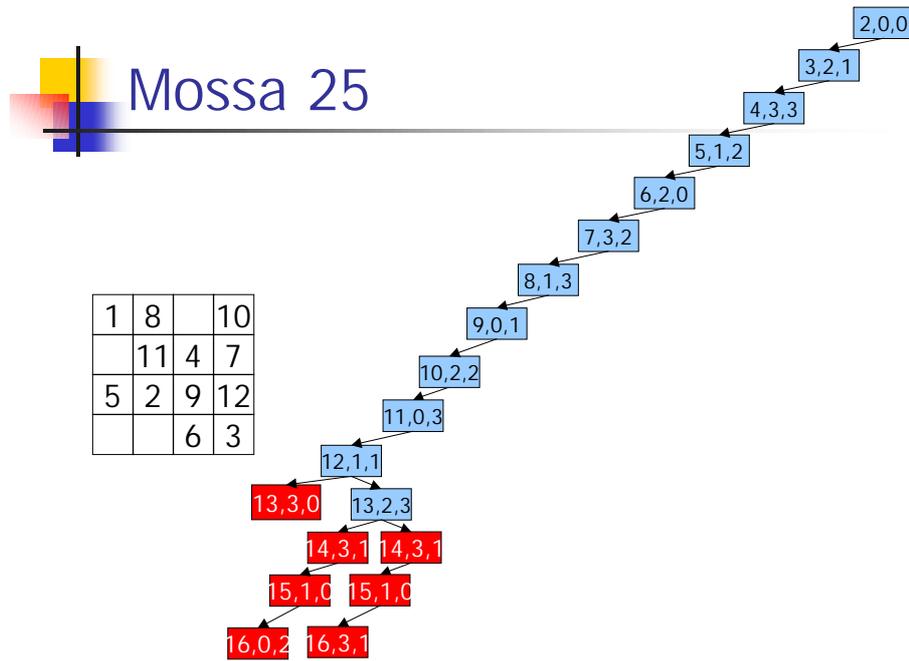
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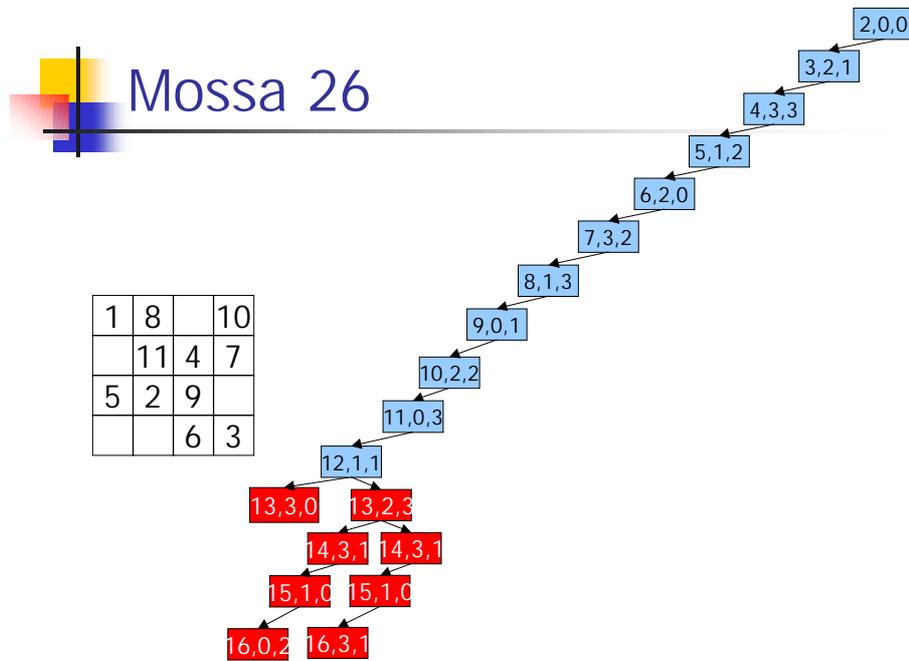
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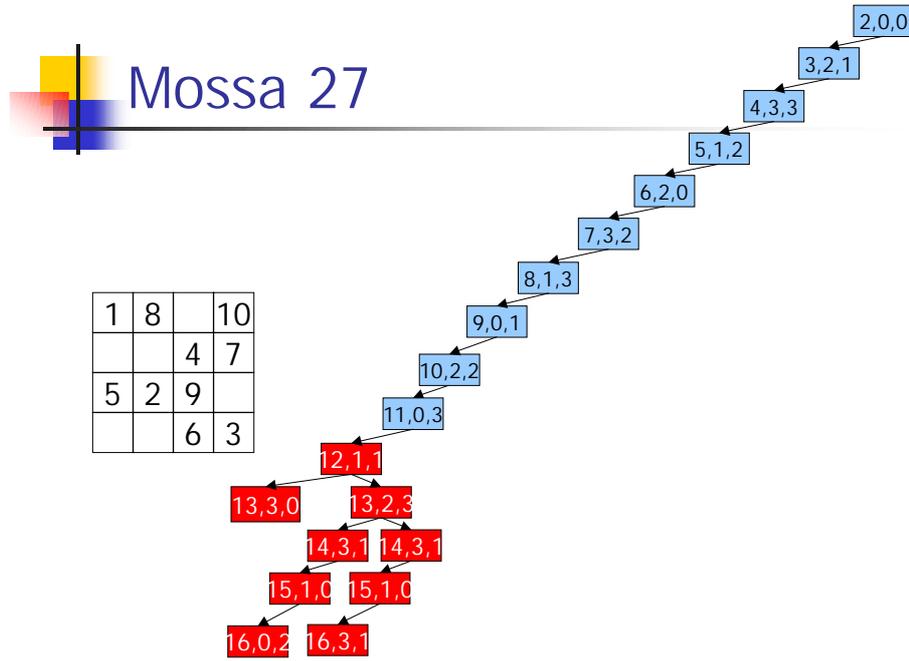
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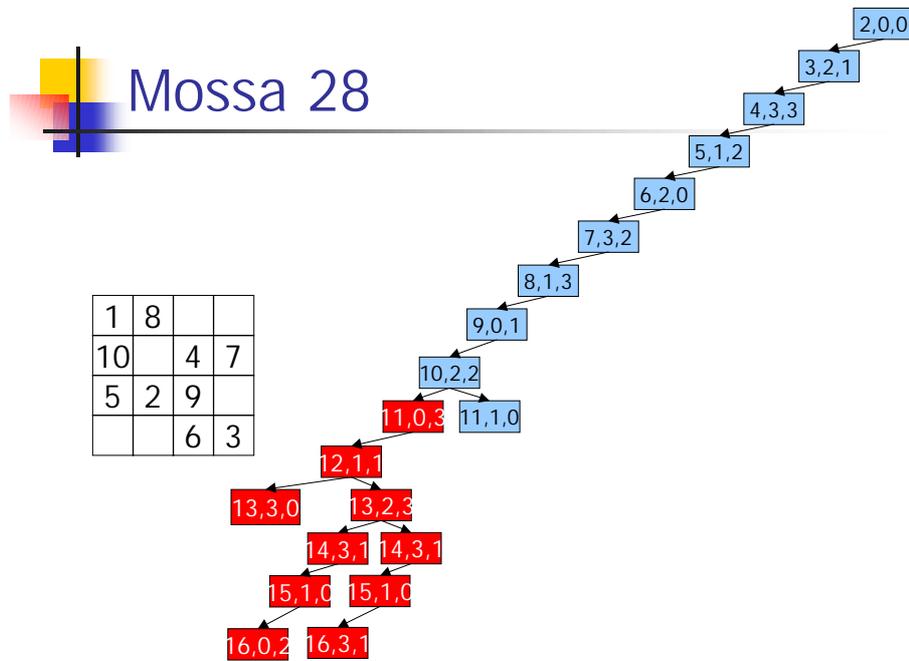
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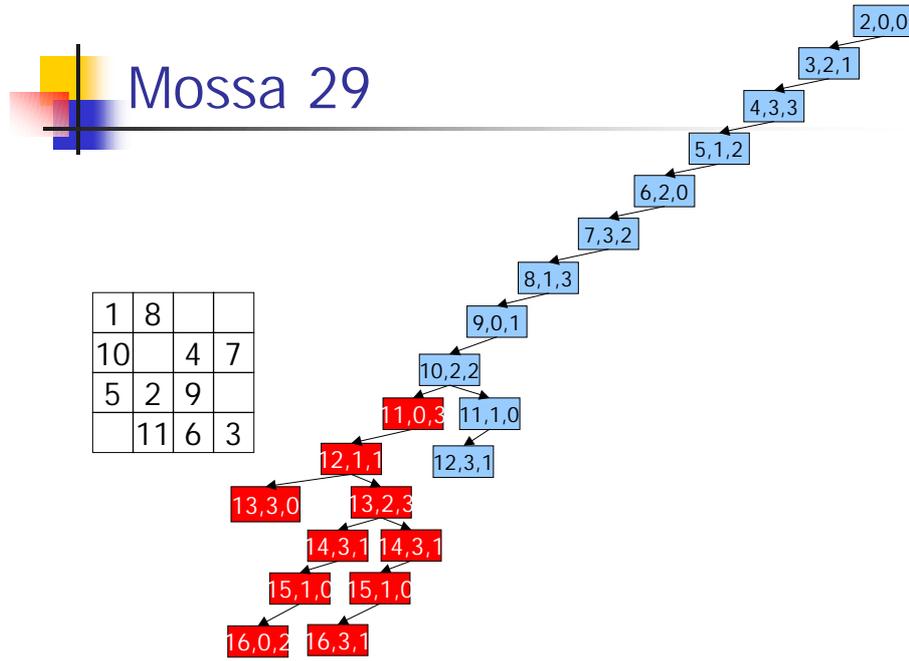
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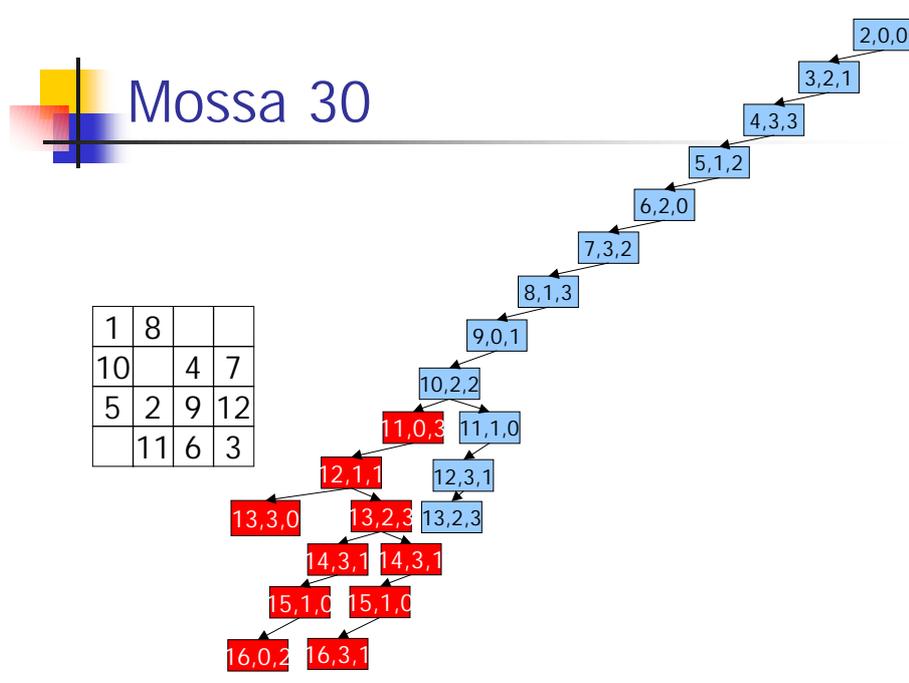
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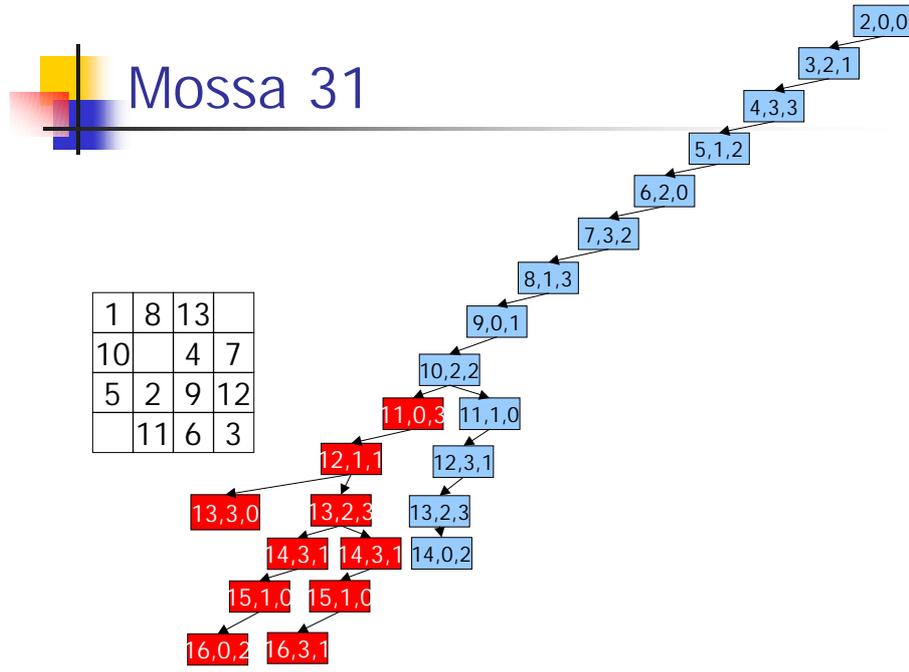
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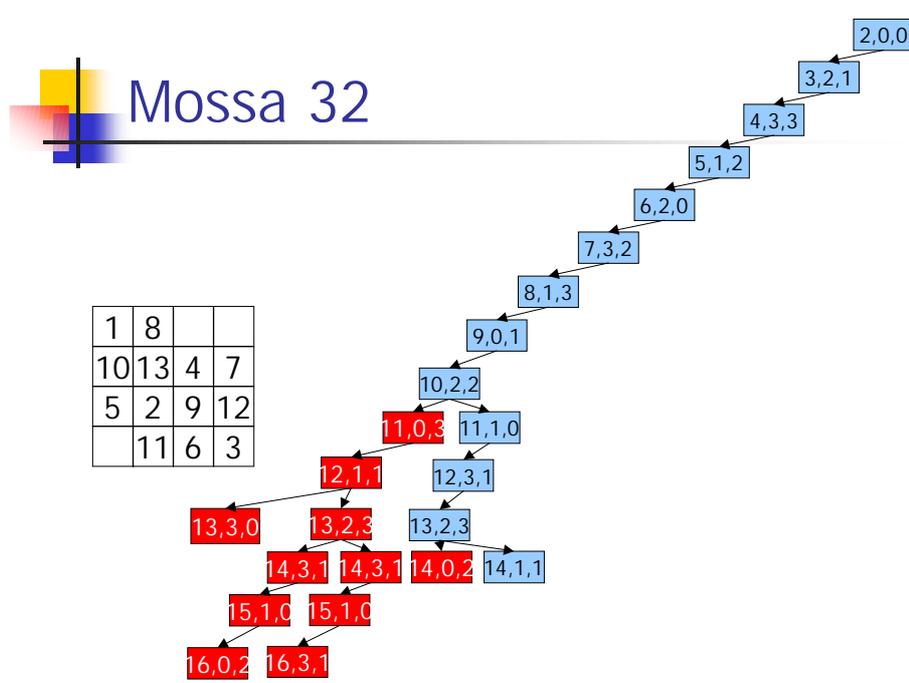
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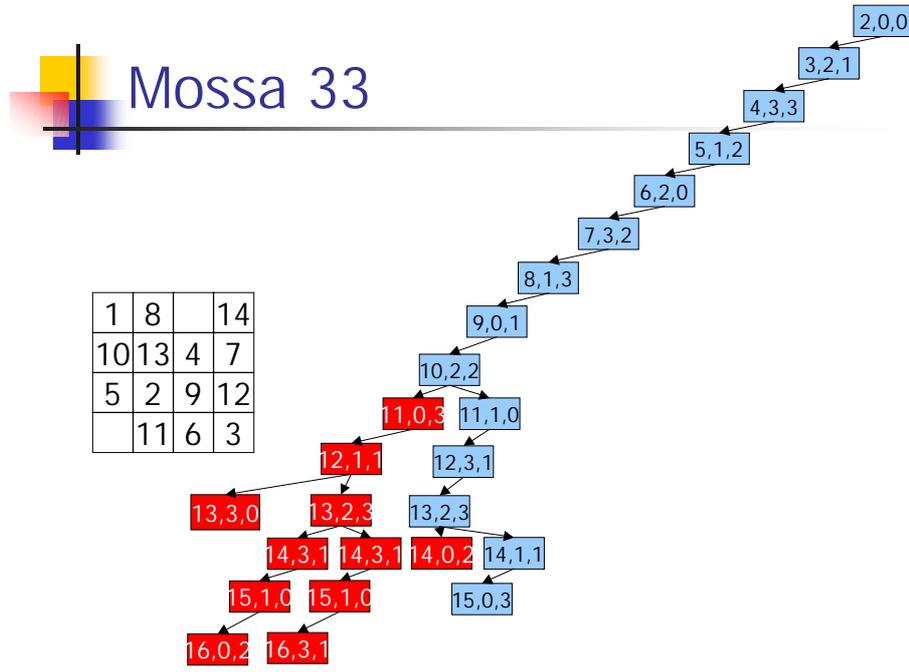
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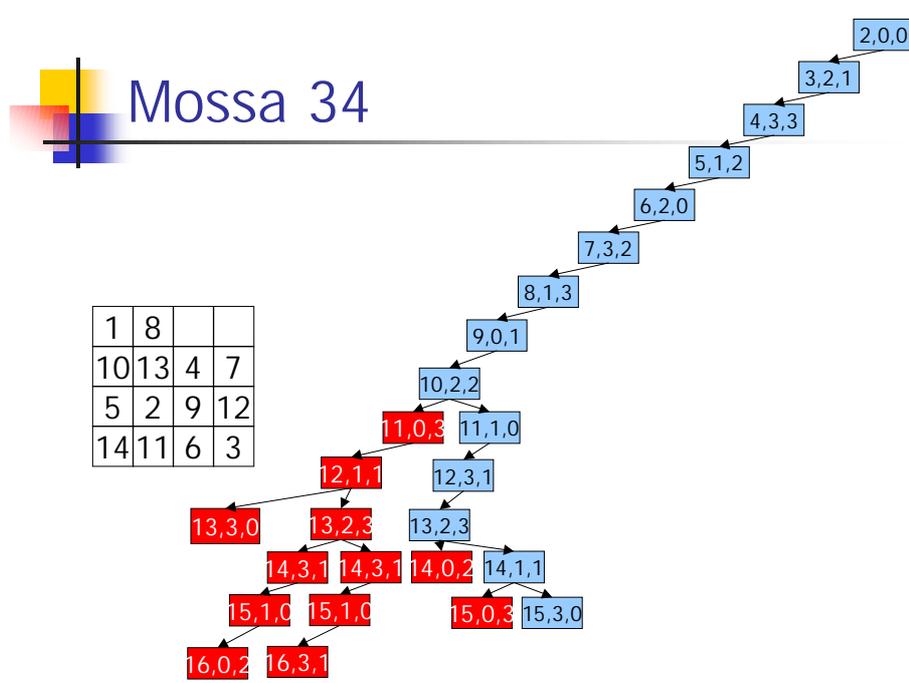
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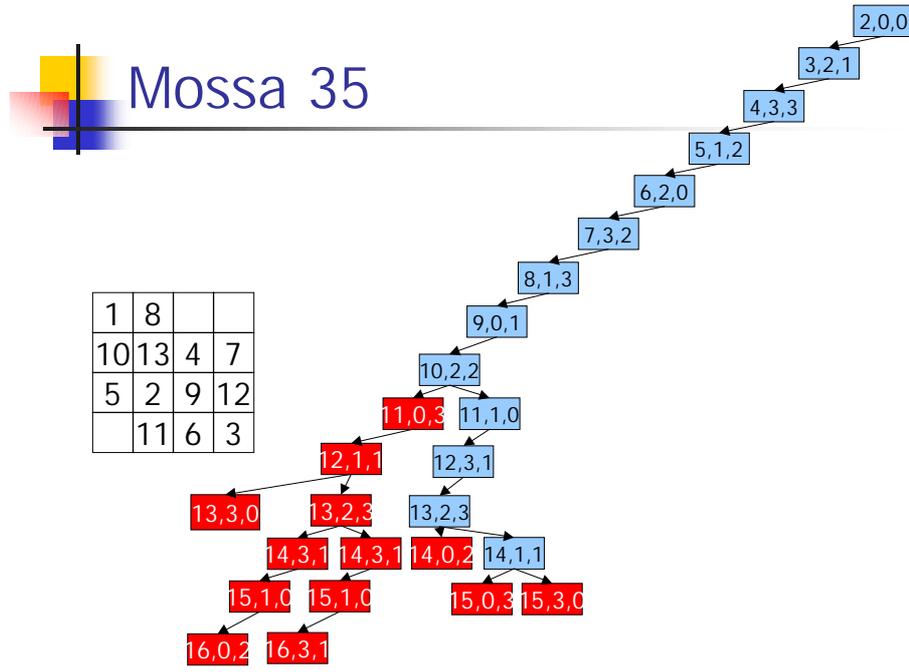
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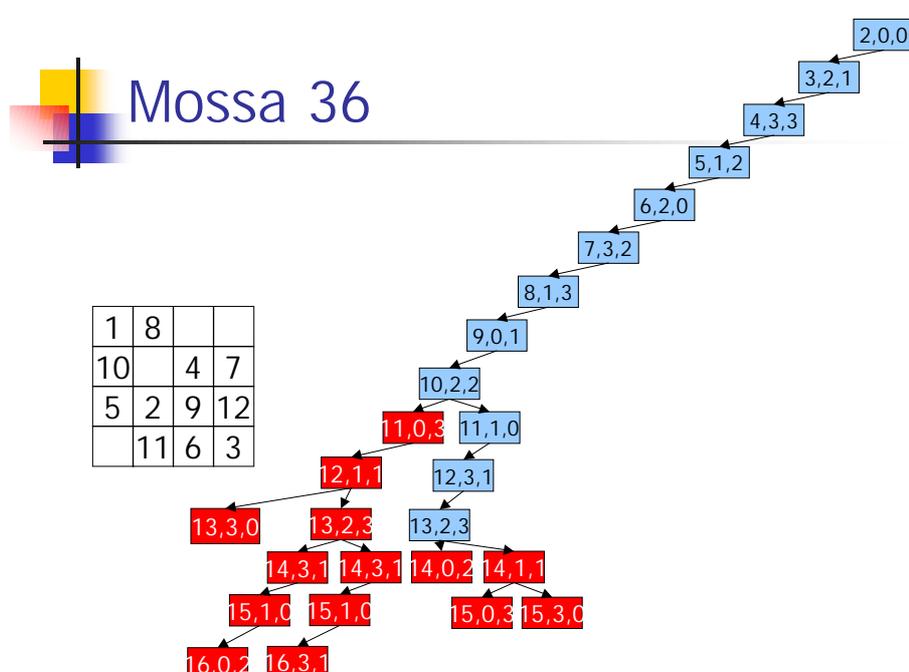
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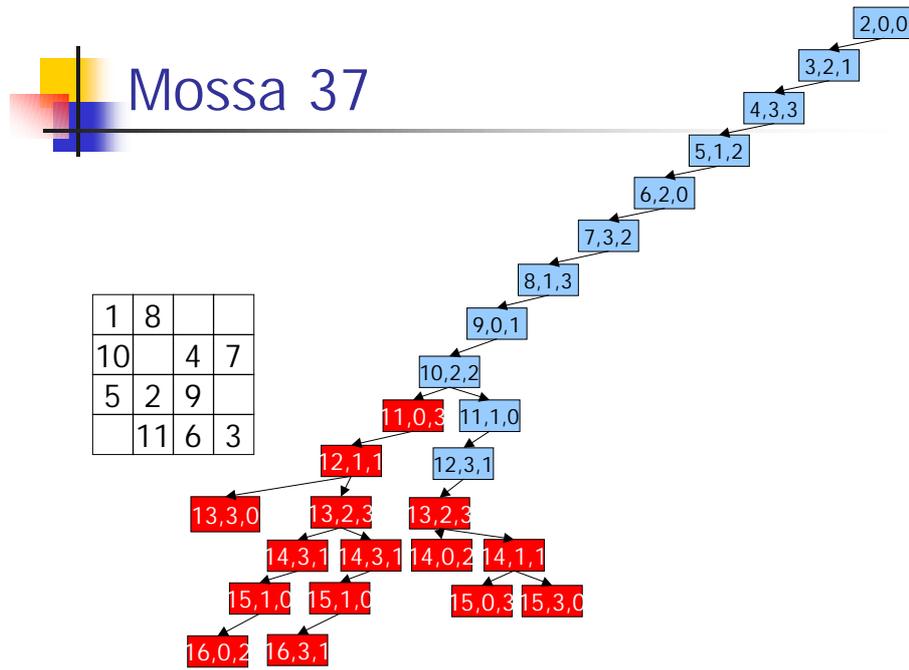
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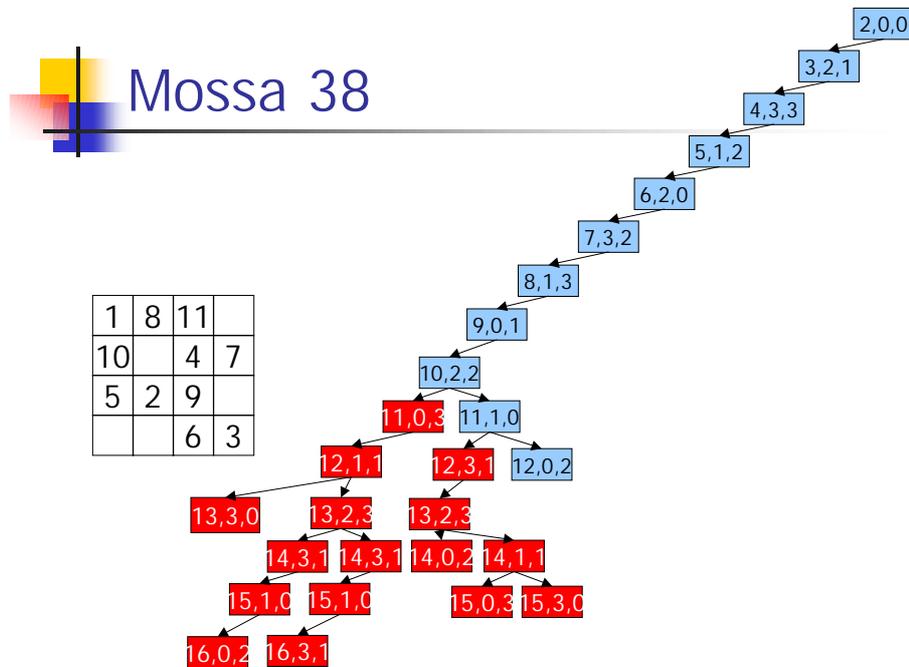
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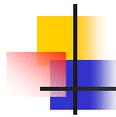
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## Complessità

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- Il numero di mosse possibili ad ogni passo è al più 8.
- Il numero di passi è  $N^2$ .
- Quindi l'albero ha un numero di nodi  $\leq 8^{N^2}$ .
- Nel caso peggiore
  - la soluzione corrisponde alla foglia più a destra
  - l'albero è completo.
- In tal caso il numero di chiamate recursive prima di trovare la soluzione è  $\Theta(8^{N^2})$ .